

F. anal. 3.4.2025

Properties of DFT?

Ex. 1.1.1?

DFT

1.

- slides, info - nests are

A. Fast Fourier transform (FFT)

DFT of $F: \mathbb{Z}(N) \rightarrow \mathbb{C}$:

$$(1) \quad F(k) = \sum_{n=0}^{N-1} c_n^N \omega_N^{-kn}, \quad c_n^N = \frac{1}{N} \sum_{k=0}^{N-1} F(k) \omega_N^{kn}, \quad \omega_N = e^{-\frac{2\pi i}{N}}$$

How to compute DFT, i.e. c_0^N, \dots, c_{N-1}^N , efficiently?

(= with fewest operations, (complex additions and multiplications))

Naive approach is $O(N^2)$:

Given $F(0), \dots, F(N-1), \omega_N$; computing

(i) ω_N^{-kn} require $N-2$ mult.

(ii) each c_n^N require $N+1$ mult., $N-1$ additions

$$\text{Total for } c_0^N, \dots, c_{N-1}^N: N-2 + N(N+1+N-1) = 2N^2 + N - 2 = O(N^2)$$

FFT (Cooley-Tukey (Gauss)) is $O(N \ln_2 N)$:

Take $N = 2^n$

$$\text{Write } \text{DFT}(N) = \frac{1}{2} (\text{DFT}(\frac{N}{2}) + \text{DFT}(\frac{N}{2}) \omega_N)$$

Solve recursively

Let $F: \mathbb{Z}(2M) \rightarrow \mathbb{C}$ and def. $F_e, F_o: \mathbb{Z}(M) \rightarrow \mathbb{C}$ by

(2) $F_e(k) = F(2k)$, $F_o(k) = F(2k+1)$

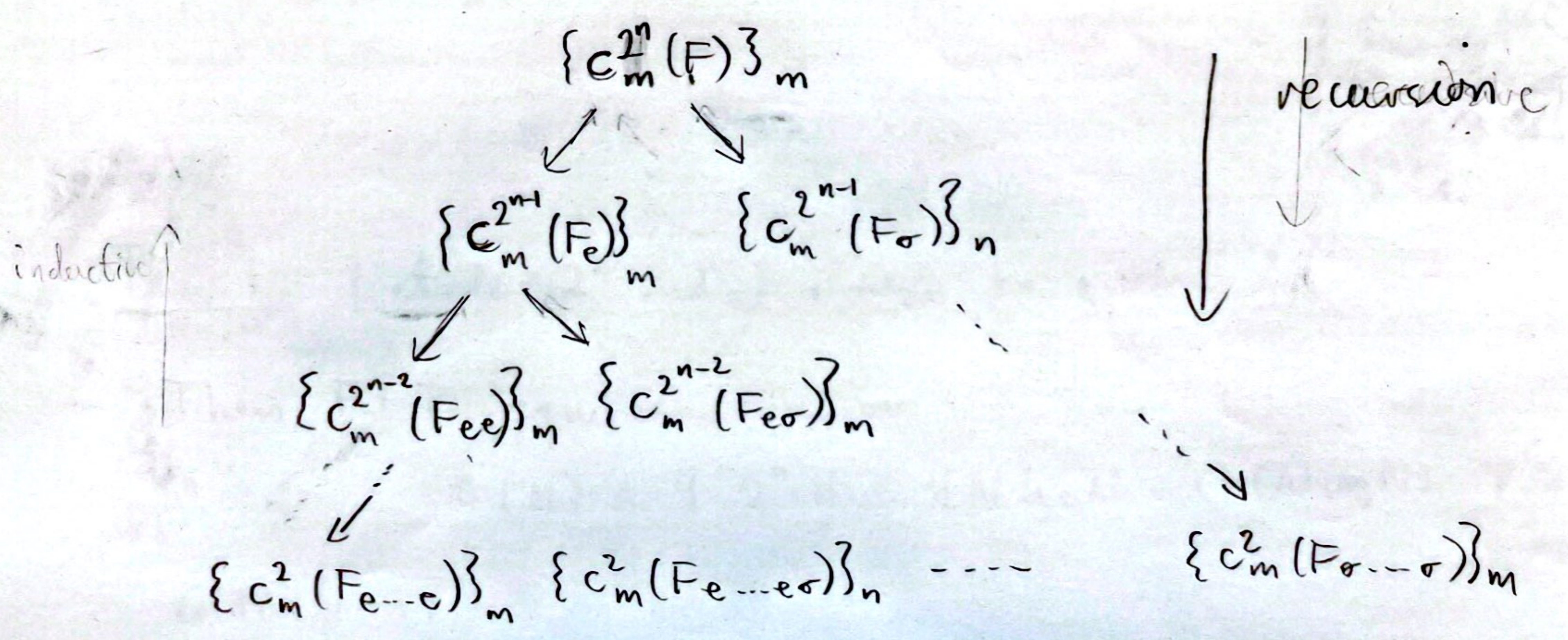
Key observation:

$$\begin{aligned}
 c_m^{2M}(F) &= \frac{1}{2M} \sum_{k=0}^{2M-1} F(k) \omega_{2M}^{km} \\
 &= \frac{1}{2} \left[\frac{1}{M} \sum_{l=0}^{M-1} F(2l) \omega_{2M}^{(2l)m} + \frac{1}{M} \sum_{\tilde{l}=0}^{M-1} F(2\tilde{l}+1) \omega_{2M}^{(2\tilde{l}+1)m} \right] \\
 &= \frac{1}{2} \left[\frac{1}{M} \sum_{l=0}^{M-1} F_e(l) \omega_M^{lm} + \frac{1}{M} \sum_{\tilde{l}=0}^{M-1} F_o(\tilde{l}) \omega_M^{\tilde{l}m} \cdot \omega_{2M} \right]
 \end{aligned}$$

(3) $= \frac{1}{2} [c_m^M(F_e) + c_m^M(F_o) \omega_{2M}]$

FFT algorithm:

Computing $c_m^N(F)$ for $N = 2^n$:



Obs: $c_m^2(G) = \frac{1}{2} (G(0) \omega_2^{0 \cdot m} + G(1) \omega_2^{1 \cdot m})$
 $= \frac{1}{2} (G(0) + (-1)^m G(1))$, $m = 0, 1$

(4)

Operations count:

#(M) = no. of operations to compute DFT on $\mathbb{Z}(M)$

Lem. 1: Given F and $\omega_{2M} = e^{-\frac{2\pi i}{2M}}$, then

$$\#(2M) \leq 2 \cdot \#(M) + 8M$$

Pf.:

(i) computing $\omega_{2M}, \dots, \omega_{2M}^{2M-1} \leq 2M$ mult's

(ii) computing $\{C_m^M(F_e)\}_m, \{C_m^M(F_o)\}_m \leq 2 \cdot \#(M)$ op's

(iii) computing $2 \cdot \#(M)$ from (3) = $2 \cdot \#(M)$ op's + 2 mult's

(iv) computing $\{C_m^{2M}(F)\}_m$ from (3)

$$= 2 \cdot \#(M) + 2M \cdot (2 \text{ mult's} + 1 \text{ addition}) + 2M$$

(ii) (3) (0)

$$= 2 \cdot \#(M) + 8M \quad \square$$

Thm. 1: Let $N = 2^n$ and F, ω_{2M} be given.

Then FFT requires $\leq 4 \cdot 2^n \cdot n = 4N \log_2 N = O(N \log_2 N)$ operations

(5) $\#(N) \leq 4 \cdot 2^n \cdot n = 4N \log_2 N = O(N \log_2 N)$ op's

Pf.:

Induction:

$n = 1$: By (4), $\#(1) \leq 2 \cdot (1 \text{ addition} + 2 \text{ mult's}) =$
 $\leq 4 \cdot 2^1 \cdot 1 = 8$ op's

so (5) holds

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Assume: (5) holds w. $N = 2^{n-1}$;

By Lem 1,

$$\#(2^n) \leq 2 \cdot \#(2^{n-1}) + 8 \cdot 2^{n-1}$$

$$\stackrel{(5) \text{ for } 2^{n-1}}{\leq} 2 \cdot (4 \cdot 2^{n-1} \cdot (n-1)) + 4 \cdot 2^n$$

$$= 4 \cdot 2^n \cdot n - 4 \cdot 2^n + 4 \cdot 2^n$$

\Rightarrow (5) holds w. $N = 2^n$

□

SS 7.3 Exercises:

① $f \in C(\mathbb{S}) \simeq C_{\text{per}}(\mathbb{R})$ (1-per. cont. f'ns on \mathbb{R})

Obs: $F(k) = f(e^{2\pi i \frac{k}{N}}) \Rightarrow F \in C(\mathbb{S}_N) \simeq C(\mathbb{Z}(N))$
 $\simeq C_{N\text{-per}}(\mathbb{Z})$

DFT of f :

$$c_n^N = \frac{1}{N} \sum_{k=0}^{N-1} f(e^{2\pi i \frac{k}{N}}) e^{-2\pi i n \frac{k}{N}}$$

F.-coeff of f :

$$c_n = \int_0^1 f(e^{2\pi i x}) e^{-2\pi i n x} dx$$

(a) $c_{n+N}^N = \frac{1}{N} \sum_k f(e^{2\pi i \frac{k}{N}}) \underbrace{e^{-2\pi i (n+N) \frac{k}{N}}}_{= e^{-2\pi i n \frac{k}{N}} \cdot \underbrace{e^{-2\pi i k}}_{= 1}} = c_n^N$

(b) $c_n^N = \sum_{k=0}^{N-1} \underbrace{f(e^{2\pi i x_k}) e^{-2\pi i n x_k}}_{\text{cont. in } x_k} \Delta x, \quad x_k = \frac{k}{N}, \Delta x = x_{k+1} - x_k = \frac{1}{N}$



$\{x_k\}_k$ partition of $[0, 1]$

$N \rightarrow \infty \Rightarrow \Delta x \rightarrow 0$

Classical results for R.-integrals

$\lim_{N \rightarrow \infty} c_n^N = \int_0^1 f(e^{2\pi i x}) e^{-2\pi i n x} dx = c_n$

② $f \in C^1(\mathbb{S})$. Show $|c_n^N| \leq \frac{C}{n}$, $0 < |n| \leq \frac{N}{2}$

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Obs: $F: \mathbb{Z}(N) \rightarrow \mathbb{C}$, then

$$\Rightarrow \sum_{k=l}^{N-1+l} F(k) = \sum_{k=0}^{N-1} F((k+l) \bmod N) = \sum_{m=0}^{N-1} F(m)$$

$$F(k \bmod N)$$

Then

$$c_n^N (1 - e^{2\pi i n \frac{l}{N}}) = \frac{1}{N} \sum_k f(e^{2\pi i \frac{k}{N}}) (e^{-2\pi i n \frac{k}{N}} - e^{-2\pi i n \frac{k-l}{N}})$$

$$\stackrel{m=k-l}{=} \frac{1}{N} \sum_k (\dots) + \frac{1}{N} \sum_{m=-l}^{N-l-l} f(e^{2\pi i \frac{m+l}{N}}) e^{-2\pi i n \frac{m}{N}}$$

$$= G: \mathbb{Z}(N) \rightarrow \mathbb{C}$$

$$\stackrel{\text{Obs}}{=} \sum_{m=0}^{N-1} f(e^{2\pi i \frac{m+l}{N}}) e^{-2\pi i n \frac{m}{N}}$$

$$\stackrel{k=m}{=} \frac{1}{N} \sum_{k=0}^{N-1} (f(e^{2\pi i \frac{k}{N}}) - f(e^{2\pi i \frac{k+l}{N}})) e^{-2\pi i n \frac{k}{N}}$$

$$\stackrel{\text{Taylor}}{=} f'(e^{2\pi i \frac{\xi}{N}}) (e^{2\pi i \frac{k+l}{N}} - e^{2\pi i \frac{k}{N}}), \xi \in (k, k+1)$$

$$= e^{2\pi i \frac{k}{N}} (e^{2\pi i \frac{l}{N}} - 1)$$

$$\stackrel{\text{Taylor}}{=} (e^x)'_{x=2\pi i \frac{\xi}{N}} \cdot \frac{2\pi i l}{N}, \xi \in (0, l)$$

Hence

$$(*) |c_n^N| \cdot |1 - e^{2\pi i n \frac{l}{N}}| \leq \frac{1}{N} \sum_{n=0}^{N-1} \|f'\|_{\infty} \cdot 1 \cdot 1 \cdot \frac{2\pi l}{N} = 2\pi \|f'\|_{\infty} \frac{l}{N}$$

For every $n, 0 < |n| \leq \frac{N}{2}$, (take $l = l(n)$ s.t.

$$\left| \frac{nl}{N} - \frac{1}{2} \right| < \frac{1}{N} \quad (\text{always ok!})$$

Then

$$\begin{aligned} |1 - e^{2\pi i \frac{nl}{N}}| &\stackrel{\Delta\text{-ineq.}}{\geq} \underbrace{|1 - e^{\pi i}|}_{=2} - \underbrace{|e^{\pi i} - e^{2\pi i \frac{nl}{N}}|}_{\stackrel{\text{Taylor}}{=} (e^z)'_{z=i\eta} \cdot (2\pi i \frac{nl}{N} - \pi)} \\ &\geq 2 - 1 \cdot 2\pi \frac{1}{N} \quad (> 0 \text{ if } N \geq 4) \end{aligned}$$

and

$$\frac{C_n}{N} \leq \frac{1}{n} \cdot \frac{1}{2} + \frac{1}{n} \left| \frac{nl}{N} - \frac{1}{2} \right| < \frac{1}{n} \left(\frac{1}{2} + \frac{1}{N} \right) \leq \frac{3}{2n}$$

Conclusion: For $N \geq 4$ and $0 < |n| < \frac{N}{2}$,

$$(*) \Rightarrow |c_n^N| \leq \frac{\frac{3}{2n}}{2 \underbrace{\left(1 - \frac{\pi}{4}\right)}_{\geq \frac{1}{2}}} \|f'\|_{\infty} \leq \frac{3\pi \|f'\|_{\infty}}{n}$$

[Optional]

(3) (a) $f \in C^2(\mathbb{S}) \Rightarrow |c_n^N| \leq \frac{C}{n^2}, 0 < |n| \leq \frac{N}{2}$

As in (2):

$$c_n^N (e^{2\pi i n \frac{L}{N}} - 2 + e^{-2\pi i n \frac{L}{N}})$$

$$= \text{as in (2)} = \frac{1}{N} \sum_{k=0}^{N-1} (f(e^{2\pi i \frac{k+L}{N}}) - 2f(e^{2\pi i \frac{k}{N}}) + f(e^{2\pi i \frac{k-L}{N}}))$$

$$|c_n^N| \cdot |e^{-2\pi i n \frac{L}{N}} - 2 + e^{2\pi i n \frac{L}{N}}|$$

Taylor

$$\leq \frac{1}{N} \sum_k \|f''\|_\infty |e^{2\pi i \frac{k}{N}} (e^{\pm 2\pi i \frac{L}{N}} - 1)|^2 \cdot |e^{-2\pi i n \frac{k}{N}}|$$

$$\leq \|f''\|_\infty \cdot 2\pi \left(\frac{L}{N}\right)^2$$

as in (2)

$$\leq \frac{C \|f''\|_\infty}{n^2}$$

(b) Inversion:

Let N be odd, and note (by Obs in (2)) that

$$(*) f(e^{2\pi i \frac{k}{N}}) = \sum_{n=0}^{N-1} c_n^N e^{2\pi i n \frac{k}{N}} = \sum_{|k| < \frac{N}{2}} c_n^N e^{2\pi i n \frac{k}{N}}, k \in \mathbb{Z}(N)$$

Obs: $c_n^N \rightarrow c_n$ by (1)

$$|c_n^N| \leq \frac{C}{n^2} \text{ by (a)}$$

Take $x_m = \frac{k(m)}{N(m)} \rightarrow x \in (0, 1)$ and $N(m) \rightarrow \infty$. By (**)

\uparrow
odd

$$f(e^{2\pi i x_m}) = \sum_{|k| < \frac{N}{2}} c_n^N e^{2\pi i x_m}$$

f cont. $\downarrow m \rightarrow \infty$

$$f(e^{2\pi i x})$$

$$\downarrow \begin{matrix} m \rightarrow \infty \\ \Rightarrow N \rightarrow \infty \end{matrix}$$

$$c_n$$

$e^{2\pi i x}$ uniformly

series abs. and unif. conv. by (a)

$$\sum_{k \in \mathbb{Z}} c_n e^{2\pi i x}$$

Hence

$$f(e^{2\pi i x}) = \sum_{k \in \mathbb{Z}} c_n e^{2\pi i x}$$